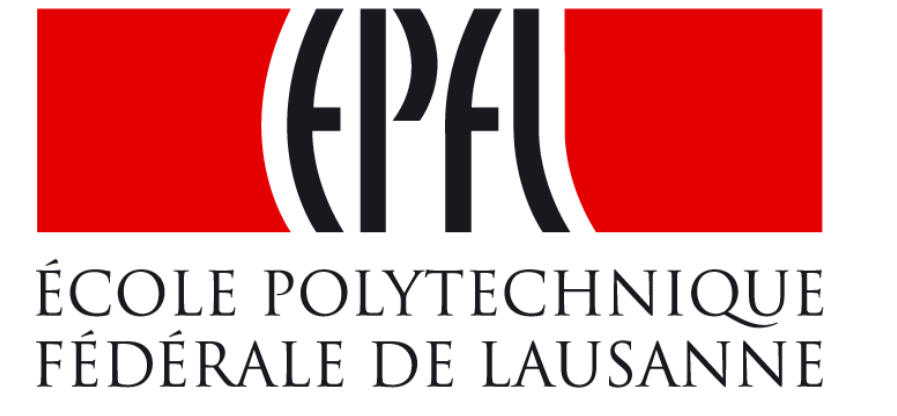


Undetectable PMU Timing-Attack on Linear State-Estimation by Using Rank-1 Approximation

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Context

- Synchrophasor measurements are vulnerable to timing attacks:

- GPS → Spoofing a GPS signal
- PTP/White Rabbit → Inserting an asymmetric delay unknown to the protocol.

△ Can we change the state of the network by just attacking the reference time of a subset of PMUs, undetectable by residual analysis?

System Model

- 1-ph direct-sequence model of a transmission network with N buses.
- Only PMU measurements (voltage or current).
- M measurements, measurement vector \mathbf{z} in \mathbb{C}^M .

Attack Model

- Attacker knows \mathbf{Y} and \mathbf{H} matrices.
- He manipulates p different time references with α_i different attack-angles ($i=1:p$).
- The attacker applies each α_i attack-angle to a subset of PMUs (\mathcal{A}_i).

- ✓ We use linear algebra with complex numbers to derive a close-form expression, and compute the attacking angles when $p=2$.

$$\alpha_1 = 2 \arg(W_{1,1} + W_{1,2}) \pmod{2\pi}$$

$$\alpha_2 = -2 \arg(W_{1,2}) + 2 \arg(W_{1,1} + W_{1,2}) \pmod{2\pi}$$

- ✓ The expression for α requires the \mathbf{W} matrix to be low rank (rank-1).

- ✓ We use the index of separation (IoS) of the \mathbf{W} matrix to derive an easy-to-do test to find vulnerable spots to attack, regardless of the state of the grid

$$IoS = \frac{\lambda_{\max}}{\sum_i \lambda_i} = \frac{\Lambda_{1,1}}{\Lambda_{1,1} + \Lambda_{2,2}} \rightarrow IoS^* = \frac{1}{2} + \frac{|f_{12}|}{2(f_{11}f_{22})^{\frac{1}{2}}} \rightarrow f_{i,j} = \sum_{l,m} \sum_n \Psi_{l,i} \Psi_{m,j} \bar{F}_{n,l} F_{n,m}$$

$$W_{i,j} = \sum_{l,m,n \in \mathcal{M}} \Psi_{l,i} \Psi_{m,j} \bar{F}_{n,l} F_{n,m} \bar{z}_l z_m$$

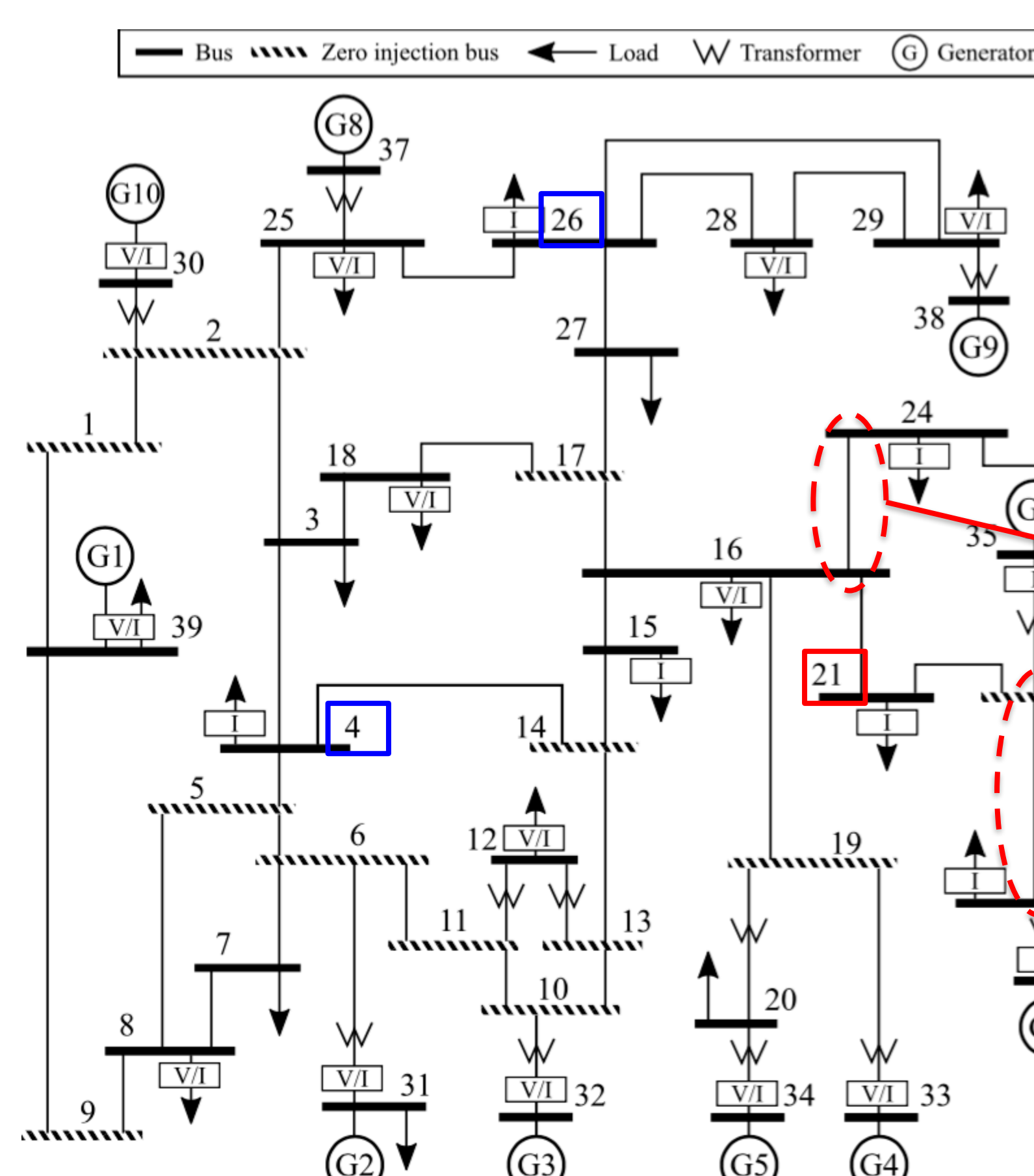
$$\Psi_{m,i} = 1 \text{ if } m \in \mathcal{A}_i \text{ and } \Psi_{m,i} = 0 \text{ otherwise}$$

$$F \triangleq H(H^H H)^{-1} H^H - I$$

Results

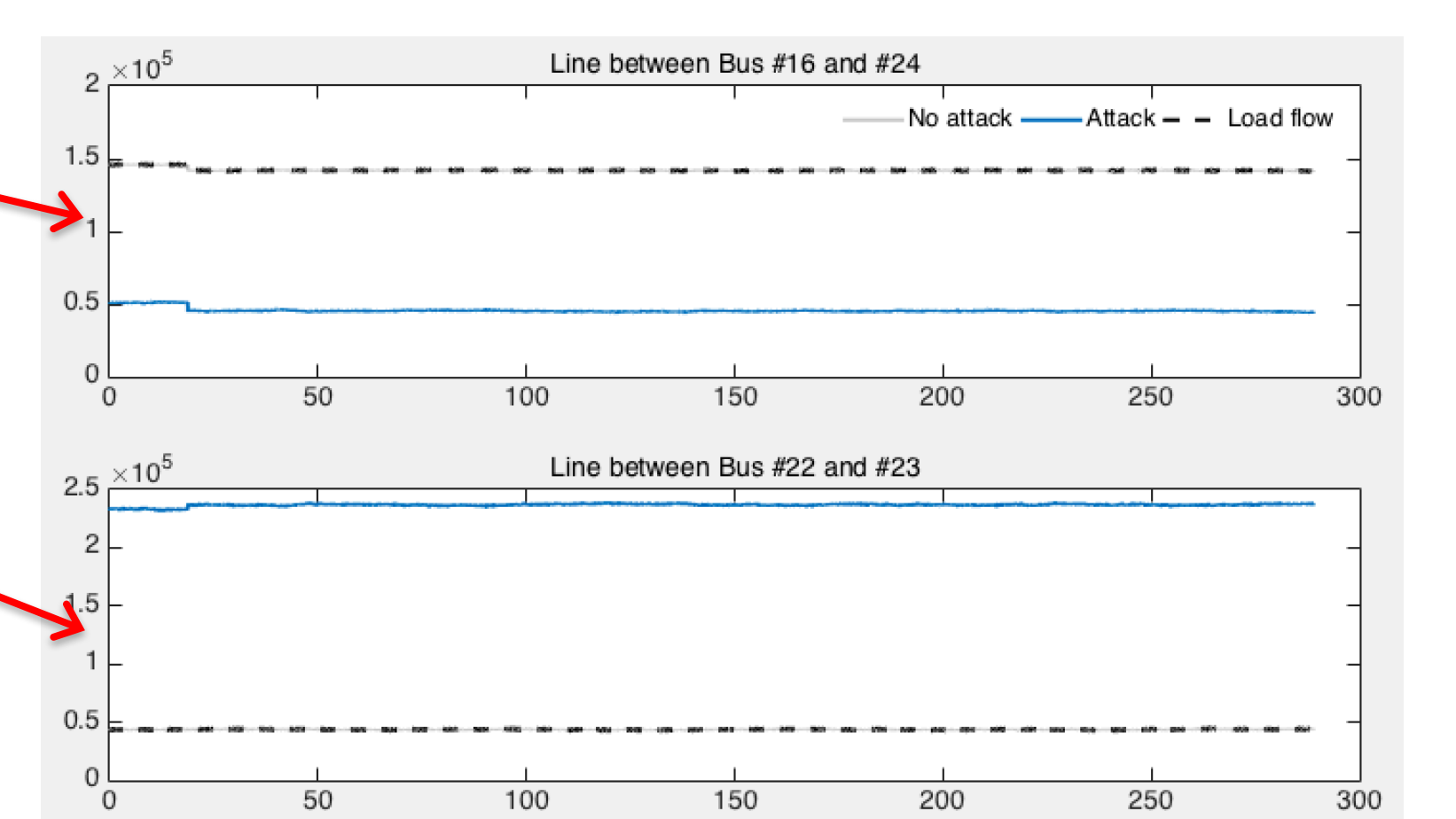
We applied IoS^* equation to all possible combination of attack locations to discover undetectable spots

Attack-location combinations					
Bus PMU1	Bus PMU2	IoS^*	Bus PMU1	Bus PMU2	IoS^*
4	15	0.8437	21	24	1.0000
4	21	0.6613	21	26	0.8395
4	23	0.6613	21	35	1.0000
4	24	0.6613	21	36	1.0000
4	26	0.5282	23	24	1.0000
4	35	0.6613	23	26	0.8395
4	36	0.6613	23	35	1.0000
15	21	0.9516	23	36	1.0000
15	23	0.9516	24	26	0.8395
15	24	0.9516	24	35	1.0000
15	26	0.7669	24	36	1.0000
15	35	0.9516	26	35	0.8395
15	36	0.9516	26	36	0.8395
21	23	1.0000	35	36	1.0000



Impact

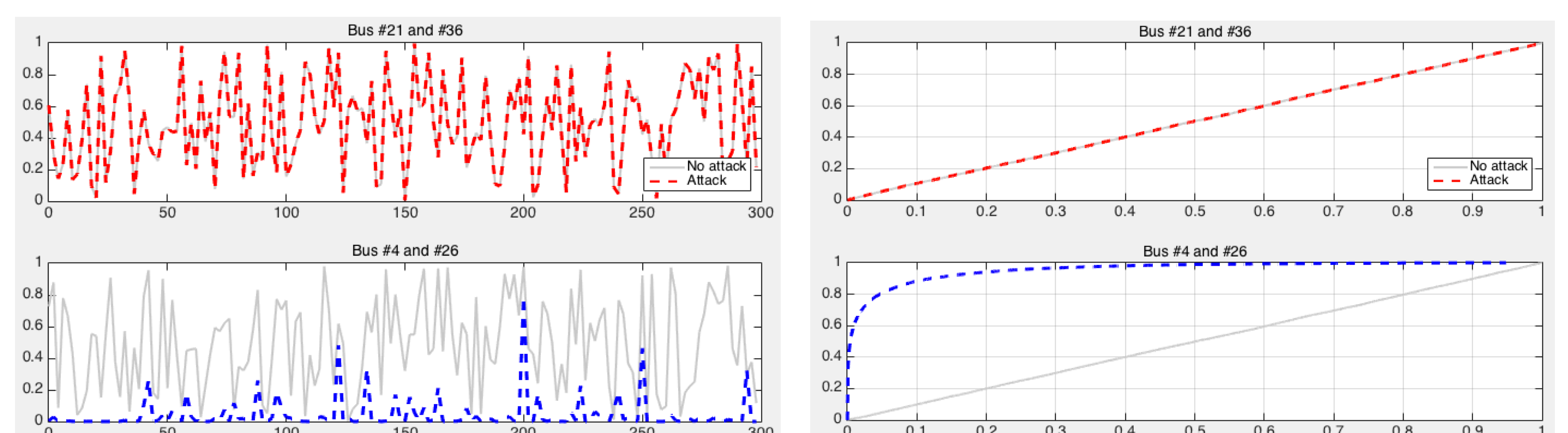
- ⚡ We deceive the network operator into believing that power flows have under- or over-utilization
- ⚡ Up to 500% error in power-flow estimation



Comparison of the true apparent power flow in two lines and the estimated apparent power flow for the no-attack and attack scenarios

Undetectability

- We use state-of-the-art bad-data detection mechanisms (i.e., χ^2 test, Largest Normalized Residual Test) to prove undetectability.
- The residuals are statistically the same before and after the attack.



Comparison of p -values and CDFs of the p -values for the χ^2 test applied to two attack locations: ideal location (in red), and lowest IoS^* -performer location (in blue). Non-attacked case is illustrated in grey.

References

- S. Barreto, M. Pignati, G. Dan, M. Paolone, J.-Y. Le Boudec, "Undetectable PMU Timing-Attack on Linear State-Estimation Using Rank-1 Approximation," submitted to IEEE Transactions on Smart Grids.
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